



# Column Generation in GAMS

## *Extending the GAMS Branch-and-Cut-and-Heuristic (BCH) Facility*

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Mathematical Optimization in Transportation  
- Airline, Public Transport, Railway -



## Welcome/Agenda

Branch-and-Cut & Heuristic Facility

First Example

Constraint Generation with BCH

Column Generation with BCH



# Agenda

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# Modeling Systems

- Best way to model and solve optimization problems
- Solid foundation based on “*Separation*”
  - Separation of Model and Data
  - Separation of Model and Algorithm
- Art of Modeling
- Some Modeling Systems provide (all) features of a programming language (e.g. GAMS, MOSEL, ...)
  - Avoid usual stumbling blocks of programming
  - Integration of optimization models
- Solver is black box
- Good approach for >95% of optimization problems
- Small number of models/users that need/want more
  - Solver/User information exchange to guide/improve the solution process.



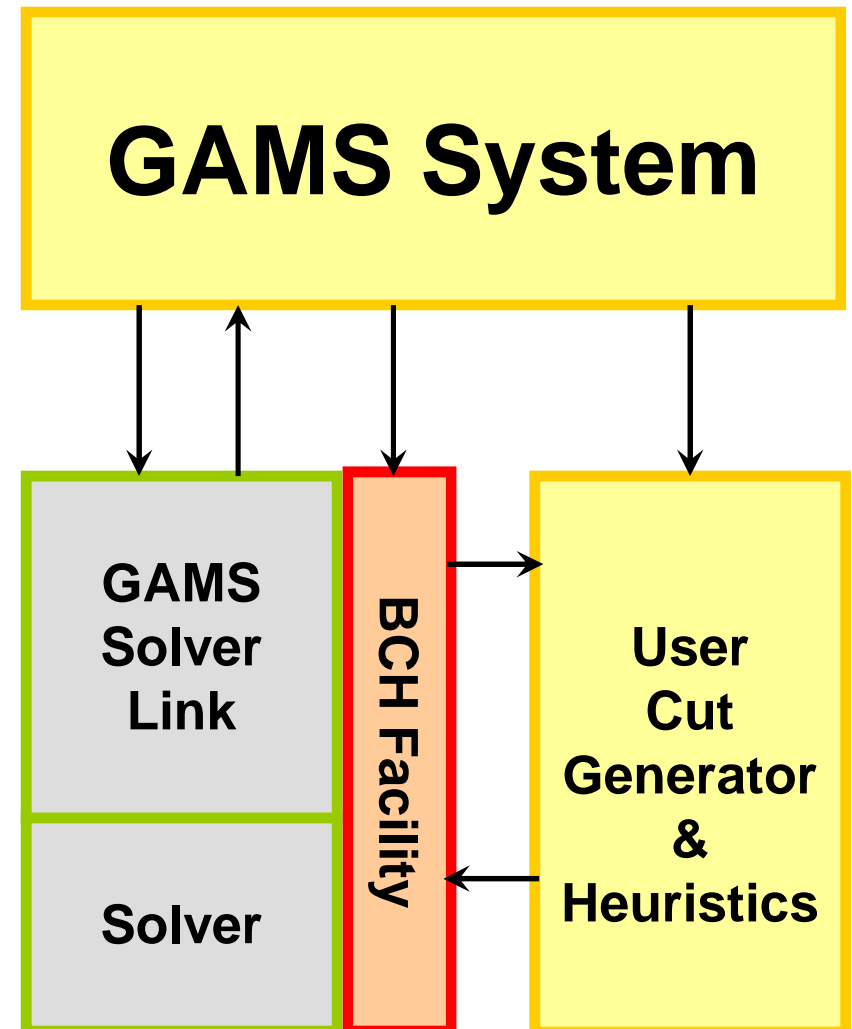
## Solution Frameworks

- Branch-and-Cut(-and-Price)
  - Abacus, MINTO
  - BCP, Bonmin, Cbc, SCIP, Symphony, ...
  - Cplex, Xpress-MP, ...
- Required Knowledge for Implementation
  - IT knowledge (C/C++/JAVA, Solver APIs)
  - Mathematical programming knowledge
  - Application specific knowledge
- Utilize rapid prototyping capability for improving solution process by user supplied information (cuts, heuristics, ...)



## “Classical” Branch-and-Cut-and-Heuristic

- Cut Generator and Heuristic
  - Represented in terms of original GAMS problem formulation
  - Independent of the specific solver
  - Use any other model type and solver available in GAMS in





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# Multi-Knapsack

<http://www.gams.com/modlib/libhtml/bchmknap.htm>

Binary variables  $x(j)$ ; Positive variables  $slack(i)$ ;  
Equations  $mk(i)$ ,  $defobj$ ; Variable  $z$ ;

```
defobj.. z =e= sum(j, value(j)*x(j));
mk(i)..  sum(j, a(i,j)*x(j)) =l= size(i);
```

```
model m /all/; solve m max z using mip;
```

**The original  
model formulation**

**Separation Problem  
for Cover Cuts:**

$$z.l < 1$$

**Cover Cuts  $c(j)=y.l(j)$ :**  
 $sum(c(j), x(j)) =l= card(j)-1$ ;

Binary variable  $y(j)$  membership in the cover;  
Equations  $defcover$ ,  $defobj$ ; Variable  $z$ ;

```
defobj..  z =e= sum(j, (1-x.l(j))*y(j));
defcover.. sum(j, ai(j)*y(j)) =g= size_i+1;
```

```
model cover /all/; solve cover min z using mip;
```





# Cover Cuts and Rounding Heuristic

- Activate BCH facility (option file):

```

usercutcall  mknap
userheurcall mknap -heuristic=1
  
```

- Separation model:

```
Excute_loadpoint 'bchout';      // Get node solution from solver
```

```
* Cover cut:
```

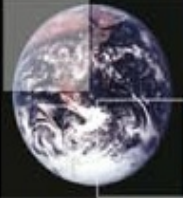
```

if (z.l<1, numcuts = 1;
  x_c('1',j) = y.l(j);                      // cut matrix
  rhs_c('1') = sum(j, y.l(j)) - 1;
  sense_c('1') = 1);                        // 1 =l=, 2 =e=, 3 =g=
  
```

```
* Heuristic
```

```

rhs(i) = rhs(i) - sum(j$(x.l(j)=1), a(i,j));
loop(j$(x.l(j)<1),
  if (smin(i, rhs(i)-a(i,j))>=0, x.l(j) = 1; rhs(i) = rhs(i) - a(i,j);
  else x.l(j) = 0));
  
```



# Cplex Log with BCH Active

	Nodes				Cuts/			
Node	Left	Objective	IInf	Best Integer	Best Node	ItCnt	Gap	
0	0	4134.0741	2		4134.0741	3		
***	Calling heuristic.	Solution obj:		3300.0000				
*	0+	0		3300.0000	4134.0741	3	25.27%	
***	Calling cut generator.	Added		2 cuts				
0	0	3871.4286	2	3300.0000	User: 1	5	17.32%	
***	Calling heuristic.	obj = 3300						
***	Calling cut generator.	Added		1 cut				
0	0	3800.0000	3	3300.0000	User: 1	7	15.15%	
***	Calling heuristic.	obj = 3300						
***	Calling cut generator.	No cuts found						
***	Calling cut generator.	No cuts found						
***	Calling heuristic.	obj = 3300						
0	2	3800.0000	3	3300.0000	3800.0000	8	15.15%	
***	Calling cut generator.	No cuts found						
***	Calling heuristic.	obj = 3800						
*	1	0	integral	0	3800.0000	3800.0000	9	0.00%

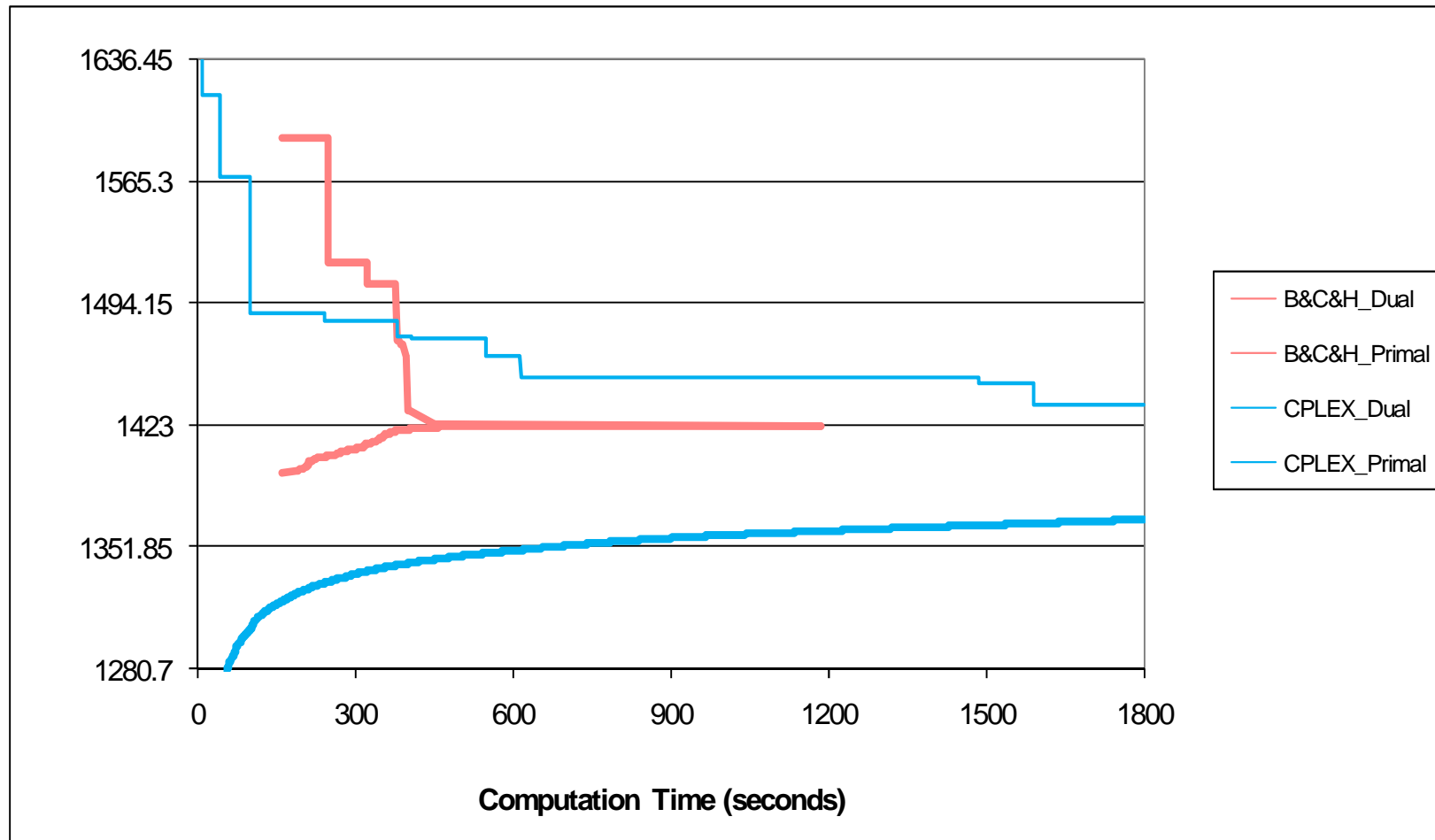


## Oil Pipeline Design Problem

- Real Example: Oil Pipeline Design Problem
  - J. Brimberg, P. Hansen, K.-W. Lih, N. Mladenovic, M. Breton 2003. An Oil Pipeline Design Problem. Operations Research, Vol 51, No. 2 228-239
  - Cuts generated when new incumbent is found
  - Rounding Heuristic, Local Branching
  - <http://www.gams.com/modlib/libhtml/bchoil.htm>
- Performance Improvements
  - Cplex/BCH: 20 minutes
  - Regular Cplex: 450 minutes
- Overhead of BCH
  - Time spent within the callback functions minus MIP computation on cuts and heuristics: 20% ~ 25%



# Oil-Design (Convergence)





## Some Recent/Ongoing Extensions

- Features
  - *Cuts and Heuristics*
  - Incumbent Filters
  - Pricing/Branching (Haase, Lübbecke, Vigerske)
  - Thread safe, BCH in a library
- Scope of Application
  - *Implement user heuristics/cuts for special problems*
  - Rapid Prototype Development for Algorithmic Ideas
    - LPEC (Michael Ferris, U Wisconsin)
    - RINS for MINLPs (Stefan Vigerske, HU Berlin)
    - Quesada/Grossmann Algorithm for MINLP
  - Constraints/Column Generation on the fly



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# Traveling Salesman Problem

- Start with “matching constraints”:

$$\sum_i x(i,j) = 1 \quad \text{for all } j$$

$$\sum_j x(i,j) = 1 \quad \text{for all } i$$

- Add subtour elimination constraints when they are needed:

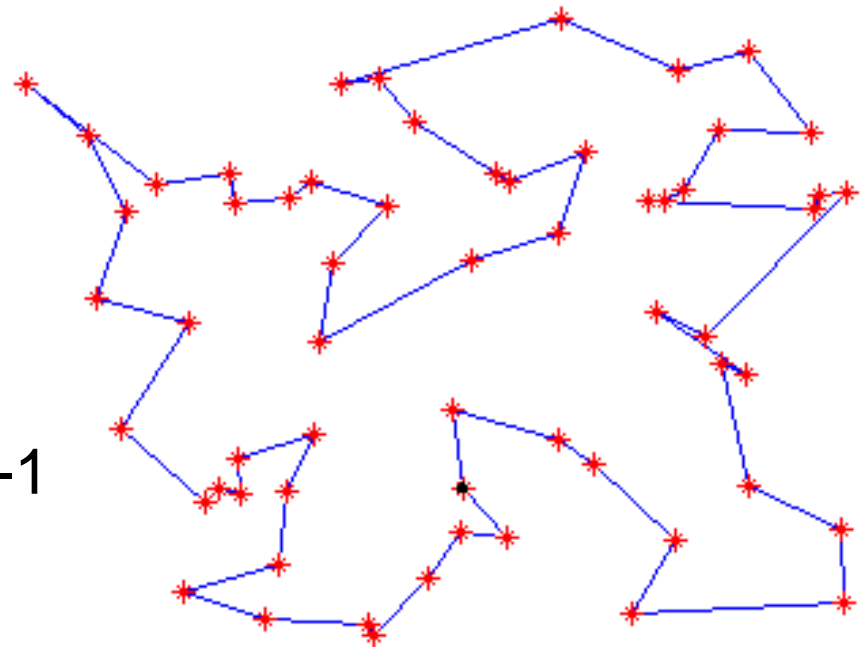
$$\sum_{i,j \in S} x(i,j) \leq \text{card}(S) - 1$$

Repeat

`solve tsp min obj using mip;`

`if (subtours, add cut);`

`until no subtours` (GAMS Model Library tsp1-tsp5)

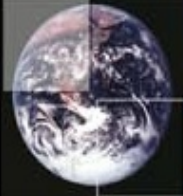




## BCH Implementation of TSP

- Perform regular B&C
  - start with matching constraints
  - presolve has to be turned off!
- Incumbent Accept/Reject Facility (`userincbcall`)
  - check for subtours
  - if rejected, store subtour elimination constraint
- Cut Facility (`usercutcall`)
  - supply subtour elimination constraint





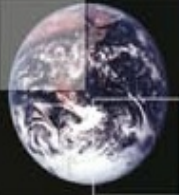
# Cplex Log for TSP

Root relaxation solution time = 0.00 sec.

\*\*\* Calling cut generator. /

\*\*\* Checking incumbent with objective 20. Rejected!

	Nodes							
	Node	Left	Objective	IInf	Best Integer	Cuts/ Best Node	ItCnt	Gap
	0	0	20.0000	0		20.0000	4	
***	Calling cut generator.		Added	3	cuts			
***	Calling cut generator.		/					
***	Checking incumbent with objective		136.	Rejected!				
***	Calling cut generator.		Added	2	cuts			
***	Calling cut generator.		/					
***	Checking incumbent with objective		78.	Accepted!				
*	0	0	integral	0	78.0000	Cuts: 13	4	0.00%
	0	0	cutoff		78.0000	78.0000	4	0.00%



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# Graph Coloring Algorithm

- Given graph  $(V,E)$ , find coloring  $c$  of  $V$  with  $c(u) \neq c(v)$  for all edges  $e=uv$  in  $E$ .
- Mehrotra/Trick:

min  $\sum(i, x(i))$ , st.  $\sum(i \text{ contains } v, x(i)) \geq 1$ ,  
 $i$  independent set of  $(V,E)$

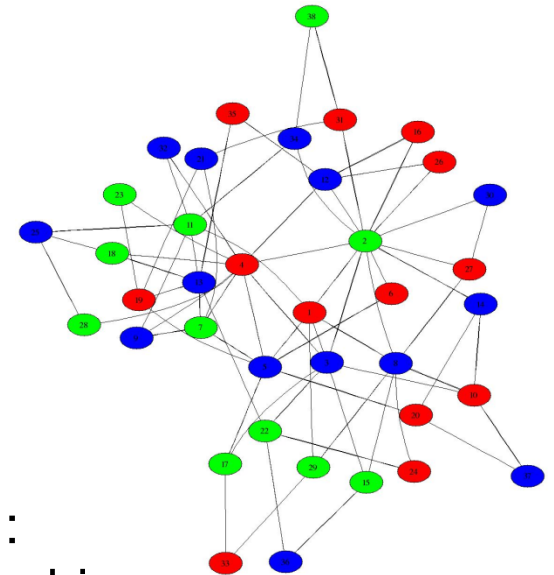
Pricing problem:

marginal cost



Find an independent set  $j$  with  $\sum(v \text{ in } j, \pi_v) > 1$ :

max  $\sum(v, \pi(v) * z(v))$ , st.  $z(u) + z(v) \leq 1$ ,  $uv$  in  $E$ ,  $z$  binary



Repeat

```

solve master min obj using rmip;
solve pricing max obj2 using mip;
if (obj2.l > 1, add column);

```

until  $\text{obj2.l} \leq 1$ ;

```

solve master min obj using mip;

```



# BCH Implementation of Graph Coloring

- Disadvantage of traditional GAMS implementation
  - Regeneration of master problem
  - Warm versus hot start of master LP
- BCH Solution:
  - Start with a trivial independent set covering
  - Solve restricted master and inside solver call pricing facility (`userpricingcall`)

Iteration	Objective	In Variable	Out Variable
1	3.252941	x91	x83
2	3.252941	defcover(20) slack	x84

```
LP status(1): optimal
Optimal solution found.
Objective :          3.252941
--- Calling pricing problem.
    3 columns added.
```

Iteration	Objective	In Variable	Out Variable
1	3.252941	x92	defcover(20) slack
2	3.250000	x80	x78

```
LP status(1): optimal
Optimal solution found.
Objective :          3.250000
```



## Outlook: Branch-and-Price

Repeat

```
  solve master min obj using rmip;  
  solve pricing max obj2 using mip;  
  if (obj2.l>1, add column);  
until obj2.l $\leq$ 1;  
solve master min obj using mip;
```

does not guarantee optimal solution



- True Branch-and-Price:
  - Column generation at the nodes of the B&B tree
  - Branching rule that is compatible with pricing problem
  - Prototype written in GAMS (tree manager, with Haase, Lübbecke)
  - Use solver framework (e.g. SCIP) to implement Branch-and-Price with BCH



## Summary

- Solver independent BCH readily available with GAMS
- Implement user heuristics, cuts and dynamic constraints/columns without too much computer science knowledge in your problem namespace
- Build rapidly prototypes of advanced algorithms in little time concentrating on the essential ideas
- Ongoing project: BCH for Branch-and-Price

<http://www.gams.com/docs/bch.htm> (documentation)

<http://www.gams.com/modlib/libhtml/alfindx.htm> (examples)